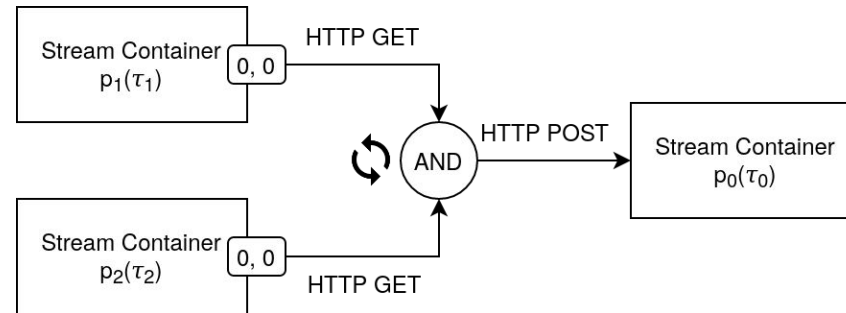


Distributed Complex Event Detection with Resource-oriented Stream Containers on the Edge



6th Stream Reasoning Workshop

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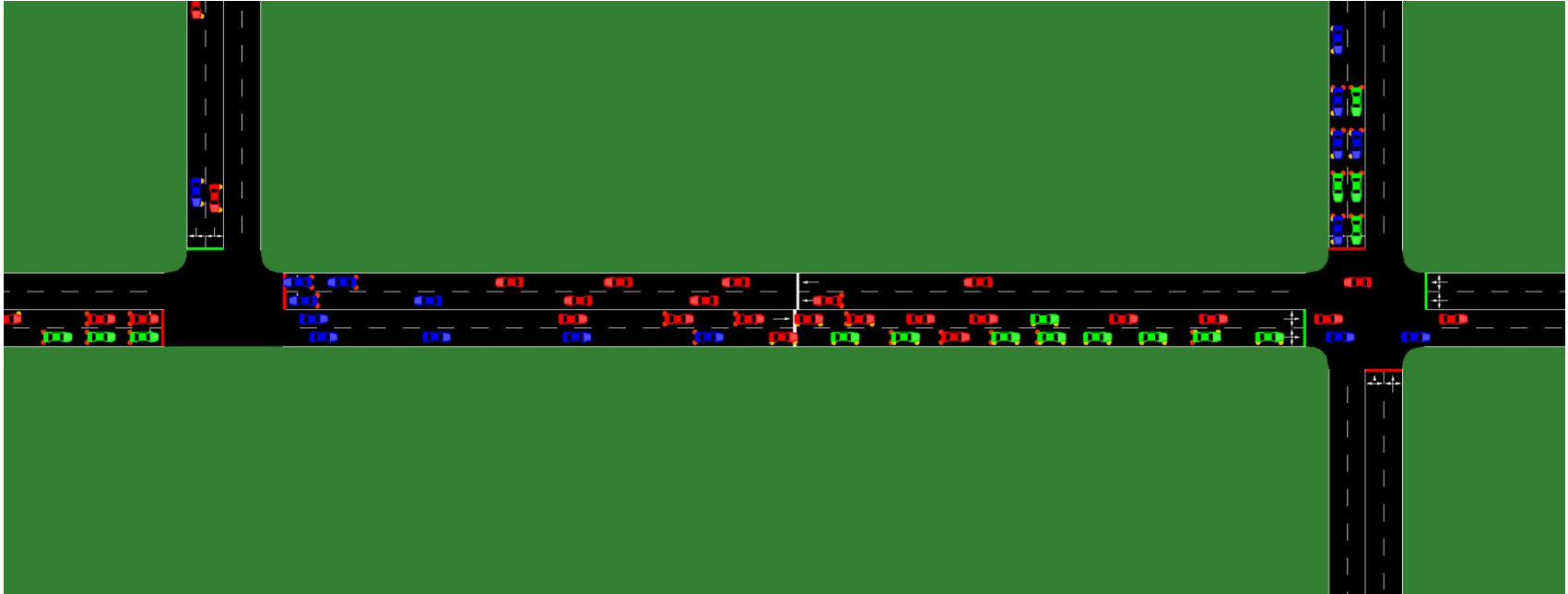
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Agenda

1. Motivation
2. Stream Containers
3. Distributed Complex Event Detection using Stream Containers
4. Splitting up MTL operators
5. Conclusion & Outlook

1. Motivation

Scenario: Traffic with many Sensors in every Car



Patrik Schneider, Daniel Alvarez-Coello, Anh Le-Tuan, Manh Nguyen Duc, Danh Le Phuoc: Stream Reasoning Playground. ESWC 2022: 406-424

Sensors & Events

Cars can have lots of sensors, e.g.

- Position
- Speed
- Acceleration
- Cardinal Direction
- Odometer
- Fuel Consumption
- Noise
- Blinker Usage
- Braking
- ...

There are lots of interesting events, e.g.

- Red-light Violations
- Turning left/right
- Stopping
- Breakdowns
- Accidents
- Speeding
- U-Turning
- Traffic Congestions
- Motor Problems
- ...

How to detect these Complex Events?

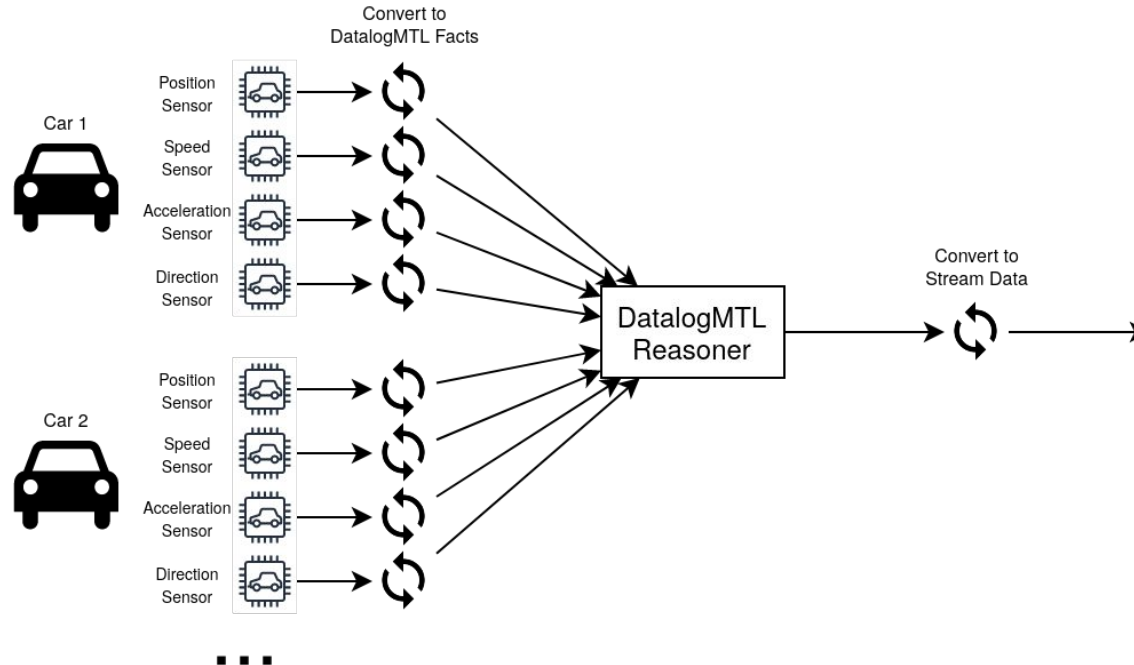
Assumption: We can describe the events using DatalogMTL^{FP} [1]

Example: Breakdown of a car

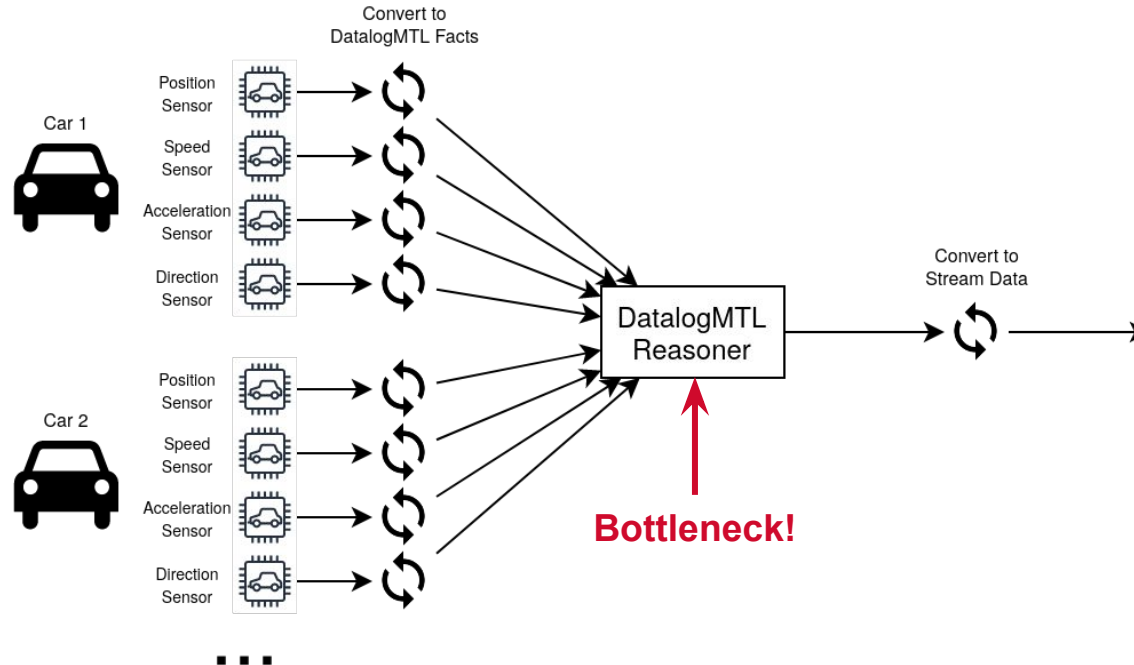
$$\textit{breakdown}(CAR) \leftarrow \exists_{[0,30s]} \textit{speed}(CAR, 0)$$

[1] Przemyslaw Andrzej Walega, Mark Kaminski, Bernardo Cuenca Grau: Reasoning over Streaming Data in Metric Temporal Datalog. AAAI 2019: 3092-3099

Streaming Complex Event: Monolithic Approach



Streaming Complex Event: Monolithic Approach



2. Stream Containers

Stream Containers [1]

- Extension of the Linked Data Platform (LDP; W3C Recommendation) in order to handle stream data in a scalable way and with a clear interface
- LDP provides a RESTful Read-Write Linked Data interface by using containers (similar to directories in a file system)
- New resources can be added to the container using HTTP POST requests

```
@prefix ldp: <http://www.w3.org/ns/ldp#> .
```

```
<>          a          ldp:BasicContainer ;  
            ldp:contains </anotherContainer>,  
                        </resource> .
```

[1] Daniel Schraudner, Andreas Harth: Stream Containers for Resource-oriented RDF Stream Processing. Stream Reasoning Workshop. 2021

Containment and Membership Triples

- Containment triple for every contained resource, membership triples depending on conditions

```
@prefix ldap: <http://www.w3.org/ns/ldap#> .  
@prefix ex: <http://example.org/> .  
@prefix sosa: <http://www.w3.org/ns/sosa/> .
```

```
<> a ldap:IndirectContainer ;  
    ldap:membershipResource <#car1> ;  
    ldap:hasMemberRelation ex:hasSpeed ;  
    ldap:insertedContentRelation sosa:hasSimpleResult ;  
    ldap:contains </obs1> , </obs2> .  
  
<#car1> ex:hasSpeed 75.4 , 77.7 .
```

```
@prefix ex: <http://example.org/> .  
@prefix sosa: <http://www.w3.org/ns/sosa/> .
```

```
</obs1> a sosa:Observation ;  
    sosa:observedProperty ex:speed ;  
    sosa:hasSimpleResult 75.4 .
```

Stream Containers use Membership Triples for Windows

```

@prefix ldapsc: <https://solid.ti.rw.fau.de/public/ns/stream-containers#> .
@prefix ldap: <http://www.w3.org/ns/ldap#> .
@prefix sosa: <http://www.w3.org/ns/sosa/> .
@prefix xsd: <http://www.w3.org/2001/XMLSchema#> .
@prefix ex: <http://example.org/> .

<>
  a          ldapsc:StreamContainer ;
  ldap:contains
    </0>,
    </1>,
    </2> ;

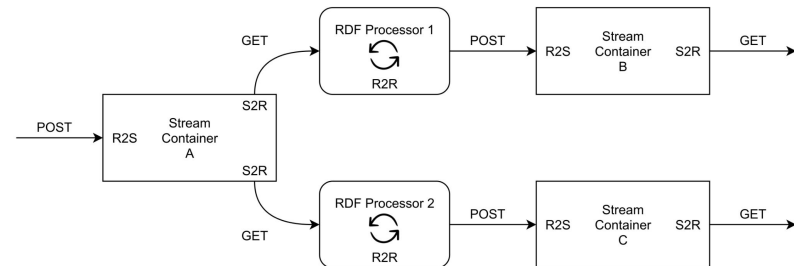
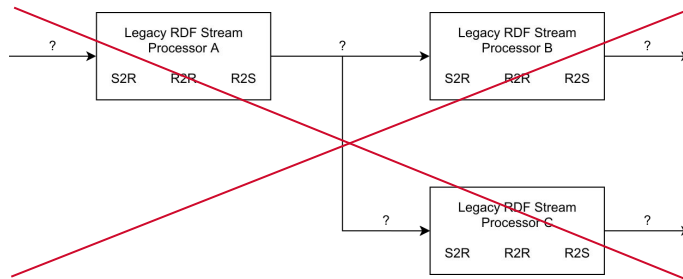
  ldapsc:window [
    ldap:hasMemberRelation      ex:inWindow ;
    ldap:membershipResource     <#window1> ;
    ldapsc:contentTimestampRelation sosa:resultTime ;
    ldapsc:beginTime            "P0S"^^xsd:duration
    ldapsc:endTime              "PT10S"^^xsd:duration
  ].

<#window1>      ex:inWindow      </1>,
                  </2> .

```

RDF Stream Processing with Stream Containers

- Stream Containers implement the R2S and S2R operators
- Together with agents capable of HTTP requests and RDF Processing (R2R) they can be used for distributed RDF Stream Processing (scalable and with clear interface)
- We were able to show that the semantics of such a system corresponds to the RSP-QL semantics



Daniel Schraudner, Andreas Harth: Stream Containers for Resource-oriented RDF Stream Processing. Stream Reasoning Workshop. 2021

3. Distributed Complex Event Detection using Stream Containers

Example: Traffic Jam

$$\text{light_jam}(CAR) \leftarrow \Diamond_{[0,15]} \text{speed_less_than_equal_30}(CAR)$$

$$\text{medium_jam}(CAR) \leftarrow \text{light_jam}(CAR) \wedge \Diamond_{[0,30]} \Box_{[0,3]} \text{speed_less_than_equal_0}(CAR)$$

$$\text{heavy_jam}(CAR) \leftarrow \text{light_jam}(CAR) \wedge \Box_{[0,30]} \Diamond_{[0,10]} \Box_{[0,3]} \text{speed_less_than_equal_0}(CAR)$$

$$\text{speed_less_than_equal_0}(CAR) \leftarrow \text{speed}(CAR, SPEED) \wedge \text{less_than_equal}(SPEED, 0)$$

$$\text{speed_less_than_equal_30}(CAR) \leftarrow \text{speed}(CAR, SPEED) \wedge \text{less_than_equal}(SPEED, 30)$$

Normal Forms

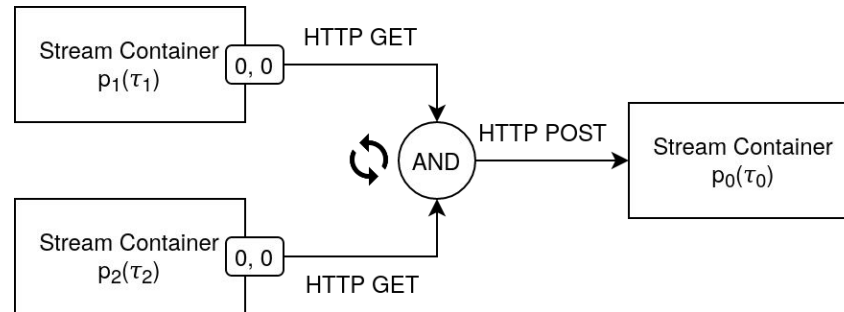
Every program in $\text{DatalogMTL}^{\text{FP}}$ with can be rewritten to a normal form. The normal form only contains rules of the form:

1. $p_0(\tau_0) \leftarrow p_1(\tau_1) \wedge \cdots \wedge p_n(\tau_n)$
2. $p_0(\tau_0) \leftarrow \exists_{\sigma} p_1(\tau_1)$
3. $p_0(\tau_0) \leftarrow \Diamond_{\sigma} p_1(\tau_1)$

Rules of Type 1

Can be handled by agents that get the current observations (window $[0,0]$) from Stream Containers and do AND

$$p_0(\tau_0) \leftarrow p_1(\tau_1) \wedge \dots \wedge p_n(\tau_n)$$

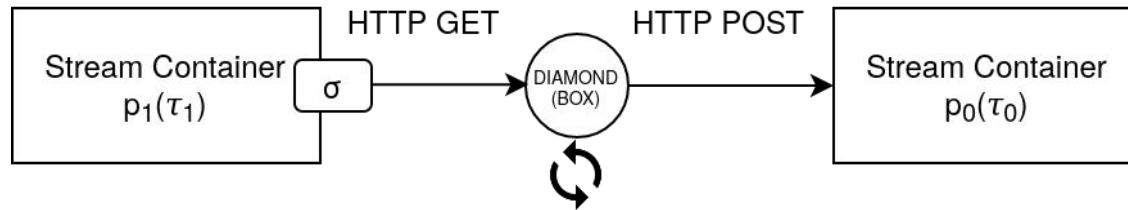


Rules of Types 2 & 3

The handling of Diamond (Box) operator from MTL is shared between Stream Containers and agents: The Stream Container calculates the window, the agent checks if there is a matching observation in the window (all observations in the window are matching).

$$p_0(\tau_0) \leftarrow \Diamond_{\sigma} p_1(\tau_1)$$

$$p_0(\tau_0) \leftarrow \Box_{\sigma} p_1(\tau_1)$$



Plan Generation

Input:

- DatalogMTL program

Output:

- List of Stream Containers with their predicates
- List of agents with their type (And, Box, Diamond), their input Stream Containers and their output Stream Container

Remarks

- The retention policy of the Stream Containers can be based on the window sizes
- Agents need a global time and a clock (e.g. every second), however, they do not need to be perfectly synchronized as they can use the HTTP Memento protocol to get the window at the correct time
- Use of background knowledge (e.g. `less_than_equal`) and OWL RL reasoning possible

Remarks

- Individual Stream Containers can be responsible for one predicate or a predicate together with a set of subjects
- It does not matter where individual Stream Containers or agents are situated - could be on the edge, could be in the cloud
- Agents are stateless and thus resilient
- We assume that this systems scales well, especially regarding the program size (different types of events to detect)

4. Splitting up MTL Operators

Splitting up the MTL Operators

If we generalize the DatalogMTL semantics like this...

$$\mathfrak{M}, t \models_{[u,v]}^{\nu} p(s, o) \text{ iff } p(\nu(s), \nu(o)) \in \mathfrak{M}(t)$$

$$\mathfrak{M}, t \models_{[u,v]}^{\nu} \Box_{[w,x]} \psi \text{ iff } \forall \min(v, t-w) \geq p \geq \max(u, t-x). \mathfrak{M}, p \models_{[\max(u, t-x), \min(v, t-w)]}^{\nu} \psi$$

$$\mathfrak{M}, t \models_{[u,v]}^{\nu} \Diamond_{[w,x]} \psi \text{ iff } \exists \min(v, t-w) \geq p \geq \max(u, t-x). \mathfrak{M} \models_{[\max(u, t-x), \min(v, t-w)]}^{\nu} \psi$$

$$\mathfrak{M}, t \models_{[u,v]}^{\nu} \Box_{[w,x]} \psi \text{ iff } \forall \max(u, t+w) \leq p \leq \min(v, t+x). \mathfrak{M} \models_{[\max(u, t+w), \min(v, t+x)]}^{\nu} \psi$$

(We get the standard semantics for $u = -\infty, v = +\infty$)

Splitting up the MTL Operators

... we can define the LTL and Window operators like this...

$$\mathfrak{M}, t \models_{[u,v]}^{\nu} W_{[w,x]}^{-} \psi \text{ iff } \mathfrak{M}, \max(u, t-x) \models_{[\max(u, t-x), \min(v, t-w)]}^{\nu} \psi$$

$$\mathfrak{M}, t \models_{[u,v]}^{\nu} W_{[w,x]}^{+} \psi \text{ iff } \mathfrak{M}, \max(u, t+w) \models_{[\max(u, t+w), \min(v, t+x)]}^{\nu} \psi$$

$$\mathfrak{M}, t \models_{[u,v]}^{\nu} \Box \psi \text{ iff } \forall t \leq p \leq v. \mathfrak{M}, p \models_{[u,v]}^{\nu} \psi$$

$$\mathfrak{M}, t \models_{[u,v]}^{\nu} \Diamond \psi \text{ iff } \exists t \leq p \leq v. \mathfrak{M}, p \models_{[u,v]}^{\nu} \psi$$

Splitting up the MTL Operators

... and show that MTL operators are equivalent to LTL + Window operators.

$$\mathfrak{M}, t \models_{[u,v]}^{\nu} \Xi_{[w,x]} \psi \iff$$

$$\forall \min(v, t - w) \geq p \geq \max(u, t - x). \mathfrak{M}, p \models_{[\max(u, t-x), \min(v, t-w)]}^{\nu} \psi \iff$$

$$\forall \max(u, t - x) \leq p \leq \min(v, t - w). \mathfrak{M}, p \models_{[\max(u, t-x), \min(v, t-w)]}^{\nu} \psi \iff$$

$$\mathfrak{M}, \max(u, t - x) \models_{[\max(u, t-x), \min(v, t-w)]}^{\nu} \Box \psi \iff$$

$$\mathfrak{M}, t \models_{[u,v]}^{\nu} W_{[w,x]}^{-} \Box \psi$$

5. Conclusion & Outlook

Conclusion

We showed how to detect complex events...

- ... based on DatalogMTL^{FP} formulas
- ... in a distributed and scalable way
- ... with a well-established and clear interface (LDP, REST)
- ... with a valid approach splitting up MTL operators in LTL and Window operators

Outlook

- Implementation ongoing
- First results of experiments regarding scalability

Thank you for your attention!